

Embedded application designed by
ATS language

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Remember Heartbleed bug?

Should we use safer language than C?

== In English ==

"Preventing heartbleed bugs with safe programming languages"

<http://bluishcoder.co.nz/2014/04/11/preventing-heartbleed-bugs-with-safe-languages.html>

== In Japanese ==

"安全なプログラミング言語を使って heartbleed を防ぐには"

<https://github.com/jats-ug/translate/blob/master/Web/bluishcoder.co.nz/2014/04/11/preventing-heartbleed-bugs-with-safe-languages.md>

"A safer systems programming language could have prevented the bug."

Want the safer language... It's ATS!

- ☆ <http://www.ats-lang.org/>
- ☆ Syntax like ML
- ☆ DML-style dependent types
- ☆ Linear types
- ☆ Optional GC
- ☆ Optional malloc/free



☆ Optional run-time

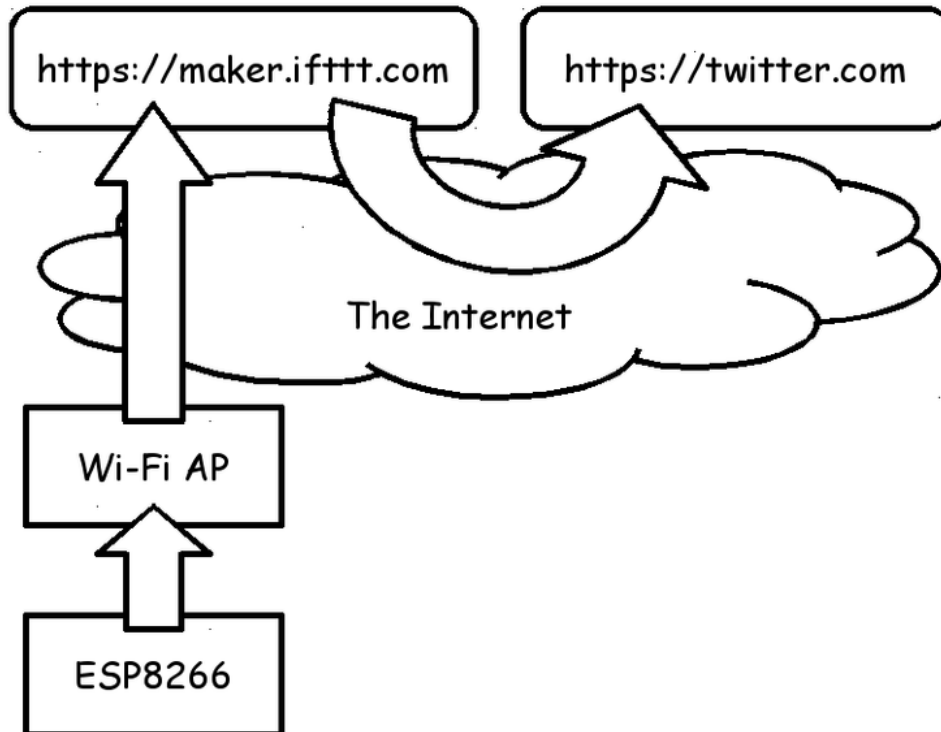
Agenda

- ☆ [1] Demo
- ☆ [2] ATS language basics
- ☆ [3] Prop (proof) on ATS
- ☆ [4] View (linear type) on ATS
- ☆ [5] Compare ATS with the other
- ☆ [6] Conclusion



[1] Demo: ATS application on ESP8266

<https://github.com/fpiot/esp8266-ats>



ESP8266 hardware

- ☆ <http://espressif.com/en/products/wroom/>
- ☆ 32-bit low power MCU Wi-Fi Module
- ☆ ROM: 4 MB
- ☆ RAM: 50 KB



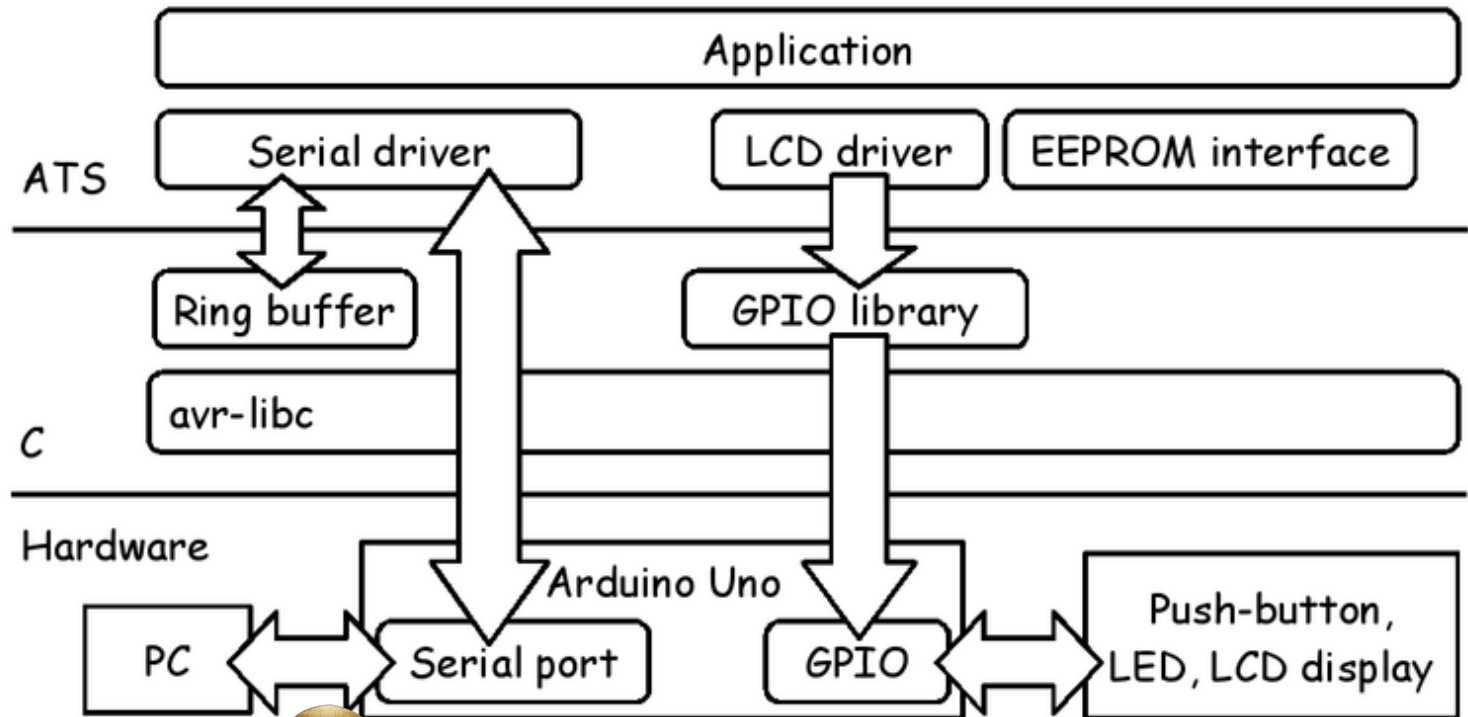
What error can be captured by ATS?

```
$ vi user/user_main.dats
39     val json_open = string0_copy "{\"value1\\": \\"
40     val json_close = string0_copy "\\ }"
41     val temp = esp_tostrptr_int rand
42     val json_head = strptr_append (json_open, temp)
43     val json_data = strptr_append (json_head, json_close)
44 //   val () = (free json_open; free json_close; free temp; free json_head)
45     val () = (free json_open; free json_close; free temp) // Error!
$ make
ATS user/user_main.dats
/home/kiwamu/src/esp8266-ats/ifttt_ats/user/user_main.dats: 985(line=32,
offs=43) -- 2237(line=55, offs=4): error(3): the linear dynamic variable
[json_head$3823(-1)] needs to be consumed but it is preserved with the type
[S2Eapp(S2Ecst(strptr_addr_vtype); S2EVar(4441))] instead.
patsopt(TRANS3): there are [1] errors in total.
```



Demo: ATS application on Arduino Uno

<https://github.com/fpiot/arduino-ats>



Arduino Uno hardware

- ☆ <https://www.arduino.cc/>
- ☆ 8-bit Harvard architecture
- ☆ ROM: 32 KB
- ☆ RAM: 2 KB

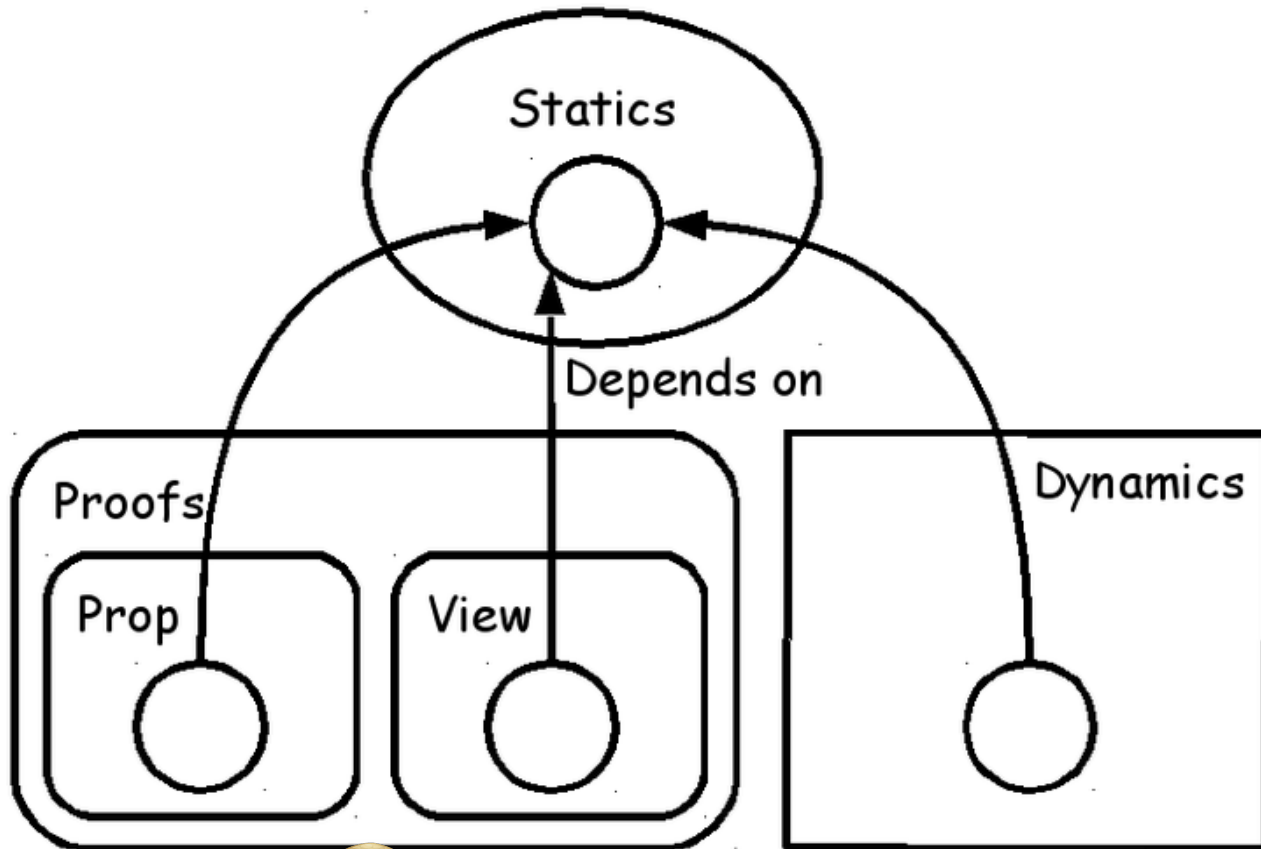


What error can be captured by ATS?

```
$ vi ../../SATS/lcd.sats
11 fun lcd_print {n:int}{i:nat | i < n}{j:nat | i + j <= n}
12 (lcd: !lcd_t, str: string (n), start: size_t (i), len: size_t (j)): void
$ vi DATS/main.dats
14 fun loop {n:int}{i:nat | i < n} .<n-i>.
15     (lcd: !lcd_t, str: string (n), pos: size_t (i)): void = {
16 // val () = if pos + i2sz LCD_WIDTH <= length str then {
17     val () = if pos + i2sz LCD_WIDTH <= 1 + length str then { // Error!
--snip--
19     val () = (lcd_setCursor (lcd, 0, 1);
20         lcd_print (lcd, str, pos, i2sz LCD_WIDTH))
$ patsopt -o DATS/main_dats.c.tmp -d DATS/main.dats
/home/kiwamu/src/arduino-ats/demo/lcd_greeting/DATS/main.dats: 1016(line=20,
offs=26) -- 1016(line=20, offs=26): error(3): unsolved constraint: C3NSTRprop
(C3TKmain()); S2Eapp(S2Ecst(<=)); S2Eapp(S2Ecst(+)); S2EVar(1969->S2Evar(i
(5685))), S2EVar(1968->S2Eintinf(16))), S2EVar(1968->S2Evar(n(5684))))
```



[2] ATS language basics



Dynamics: ML-style programming

- ☆ Dynamics of ATS is similar to Standard ML.
- ☆ You should represent type signature of function, because ATS can't inference everything.
- ☆ You should introduce main function, because ATS code is compiled into C language.



Fizzbuzz on Standard ML

local

```
fun fbstr i =  
  case (i mod 3 = 0, i mod 5 = 0) of  
    (true , true ) => "FizzBuzz"  
  | (true , false) => "Fizz"  
  | (false, true ) => "Buzz"  
  | (false, false) => Int.toString i
```

```
fun fizzbuzz' (n, j) =  
  if n = j then ()  
  else (print (fbstr j ^ "\n"); fizzbuzz' (n, j+1))
```

in

```
fun fizzbuzz n = fizzbuzz' (n, 1)  
val _ = fizzbuzz 100
```

end



Fizzbuzz on ATS

local

```
fun fbstr (i:int): string =  
  case (i mod 3 = 0, i mod 5 = 0) of  
    (true , true ) => "FizzBuzz"  
  | (true , false) => "Fizz"  
  | (false, true ) => "Buzz"  
  | (false, false) => toString_int i
```

```
fun fizzbuzz' (n:int, j:int): void =  
  if n = j then ()  
  else (println! (fbstr j); fizzbuzz' (n, j+1))
```

in

```
fun fizzbuzz (n:int): void = fizzbuzz' (n, 1)  
val _ = fizzbuzz 100
```

end



ATS binary is portable

☆ File size is 13 kB:

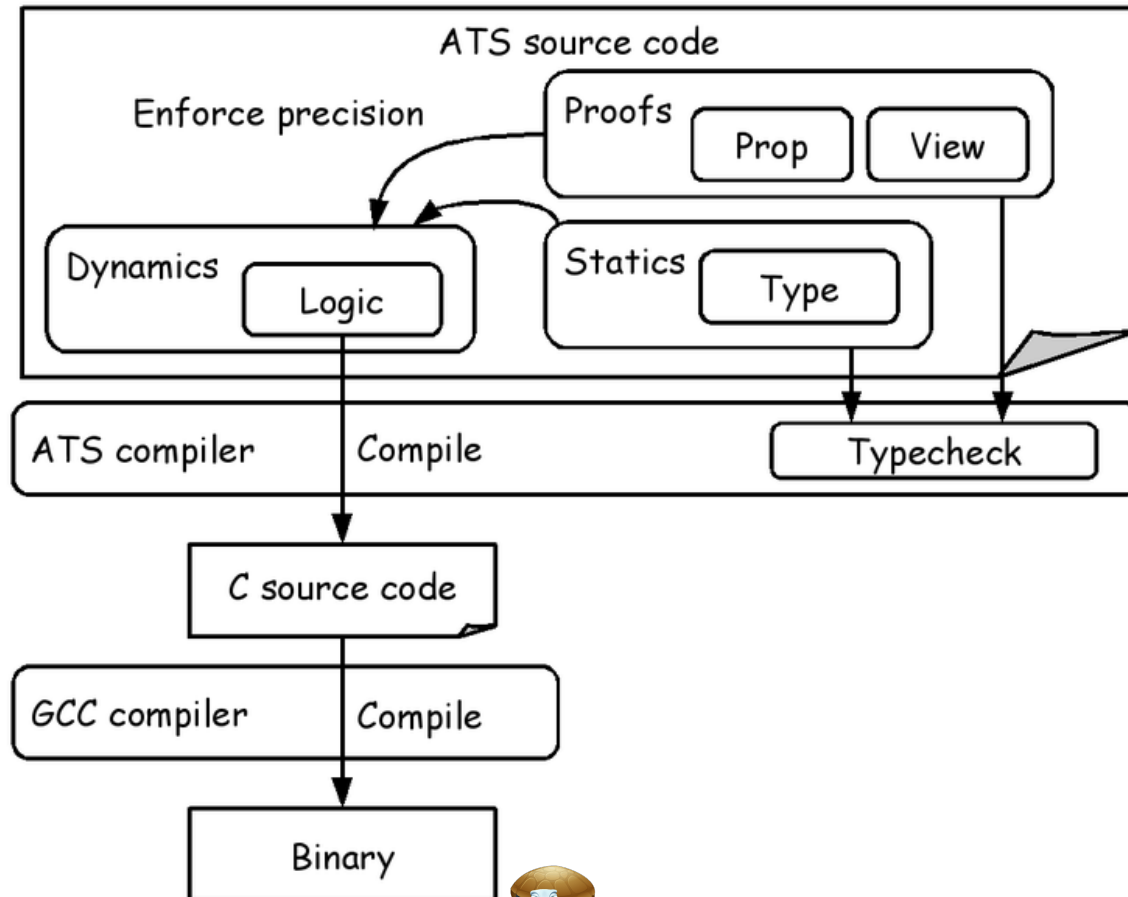
```
$ ls -lh fizzbuzz
-rwxr-xr-x 1 kiwamu kiwamu 13K Nov  6 21:45 fizzbuzz*
```

☆ Number of undefined symbols is 6:

```
$ nm fizzbuzz | grep 'U '
U __libc_start_main@@GLIBC_2.2.5
U exit@@GLIBC_2.2.5
U fflush@@GLIBC_2.2.5
U fprintf@@GLIBC_2.2.5
U fwrite@@GLIBC_2.2.5
U longjmp@@GLIBC_2.2.5
```



ATS compile flow

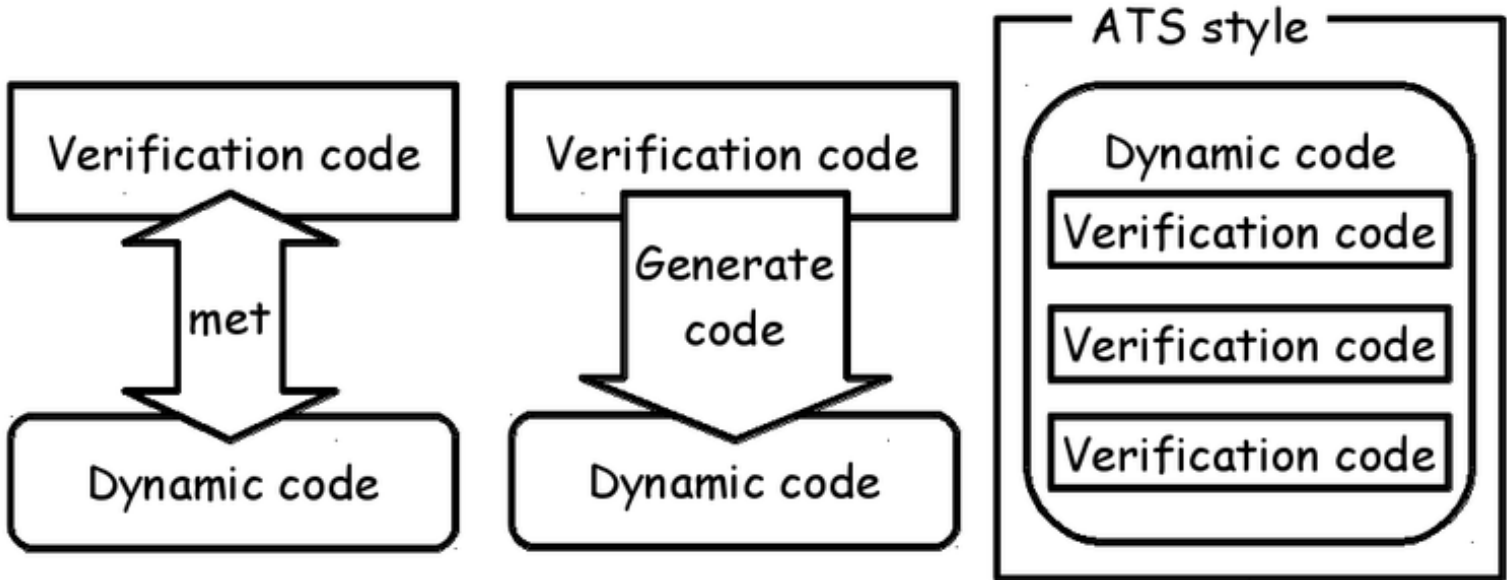


What can ATS do?

- ☆ Write code on bare metal hardware.
- ☆ Write code in Linux kernel.
- ☆ Use strong type without any OS.
- ☆ Prove code using dependent types.
- ☆ Safely use malloc using linear types.
- ☆ Safely use pointer using linear types.



[3] Prop (proof) on ATS

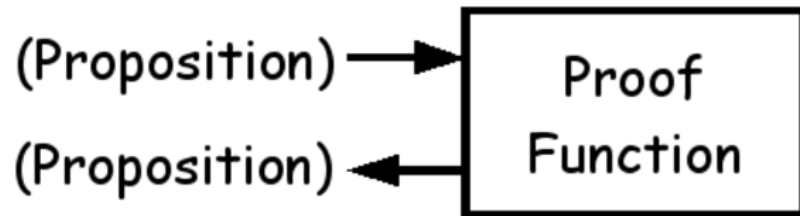
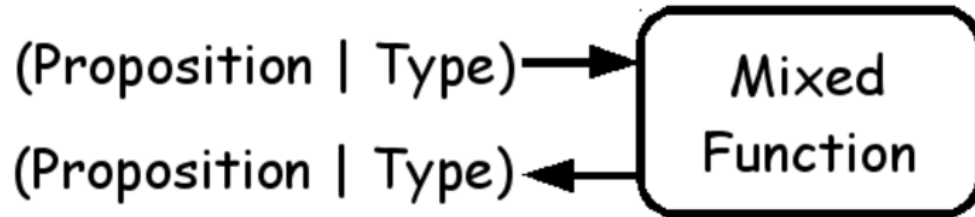
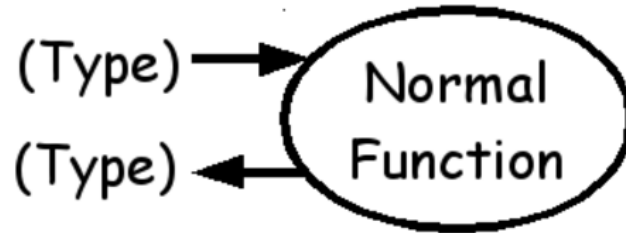


Prop: Curry–Howard in ATS

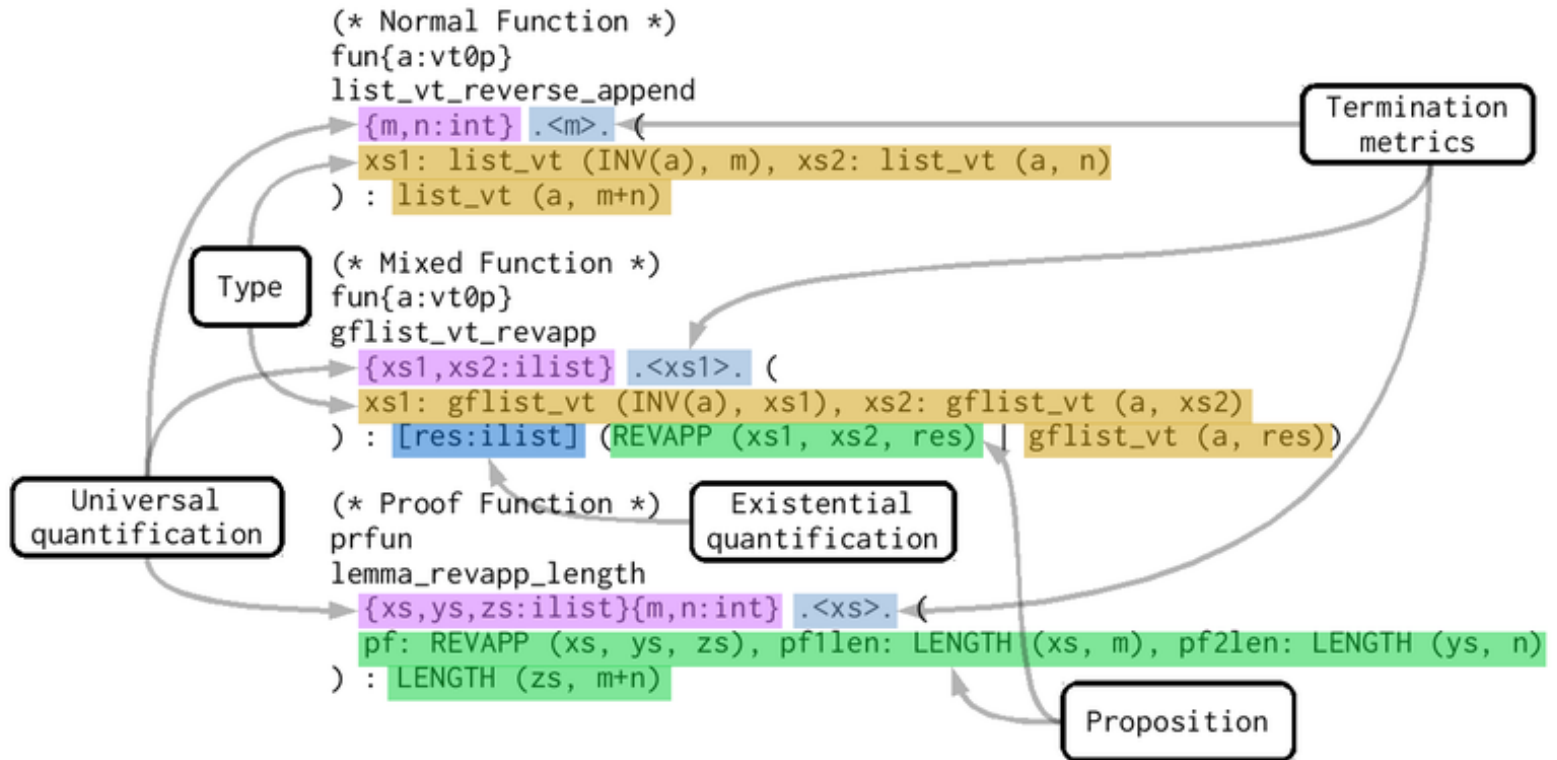
- ☆ Type: Function signature introduced by keyword "fun"
- ☆ Program: Function body introduced by keyword "implement"
- ☆ Proposition: Proof function signature introduced by keyword "prfun"
- ☆ Proof: Proof function body introduced by keyword "primplement"



Prop: Style of functions



Prop: Function signature



Prop: Function body

```
fun{a:vt0p}
split
  {xs:ilist}{n,i:nat | i <= n} .<i>. (
    pflen: LENGTH (xs, n)
  | xs: &gflist_vt (a, xs) >> gflist_vt (a, xs1), i: int i
  ) : #[xs1,xs2:ilist] (
    APPEND (xs1, xs2, xs), LENGTH (xs1, i) | gflist_vt (a, xs2)
  ) =
  if i > 0 then let
    prval LENGTHcons (pflen) = pflen
    val @gflist_vt_cons (_, xs1) = xs
    val (pfapp, pf1len | xs2) = split (pflen | xs1, i-1)
    prval () = fold@ (xs)
  in
    (APPENDcons (pfapp), LENGTHcons (pf1len) | xs2)
  end else let
    val xs2 = xs
    val () = xs := gflist_vt_nil ()
    prval pfapp = append_unit_left ()
  in
    (pfapp, LENGTHnil () | xs2)
  end
end
```

Proof

Program



Prop is similar to Coq

- ☆ Coq is an interactive theorem prover.
- ☆ <https://coq.inria.fr/>
- ☆ ATS also has subsystem for theorem-proving.
- ☆ But ATS doesn't have tactics.
- ☆ ATS constructs proofs as total functions.



Prop: Coq code

Inductive day : Type :=

- | monday : day
- | tuesday : day
- | wednesday : day
- | thursday : day
- | friday : day
- | saturday : day
- | sunday : day.

Inductive good_day : day -> Prop :=

- | gd_sat : good_day saturday
- | gd_sun : good_day sunday.

Theorem gds : good_day sunday.

Proof. apply gd_sun. Qed.



Prop: ATS code

datasort Day =

- | Monday
- | Tuesday
- | Wednesday
- | Thursday
- | Friday
- | Saturday
- | Sunday

dataprop Good_Day (Day) =

- | Gd_Sat (Saturday) of ()
- | Gd_Sun (Sunday) of ()

extern prfun gds: Good_Day Sunday

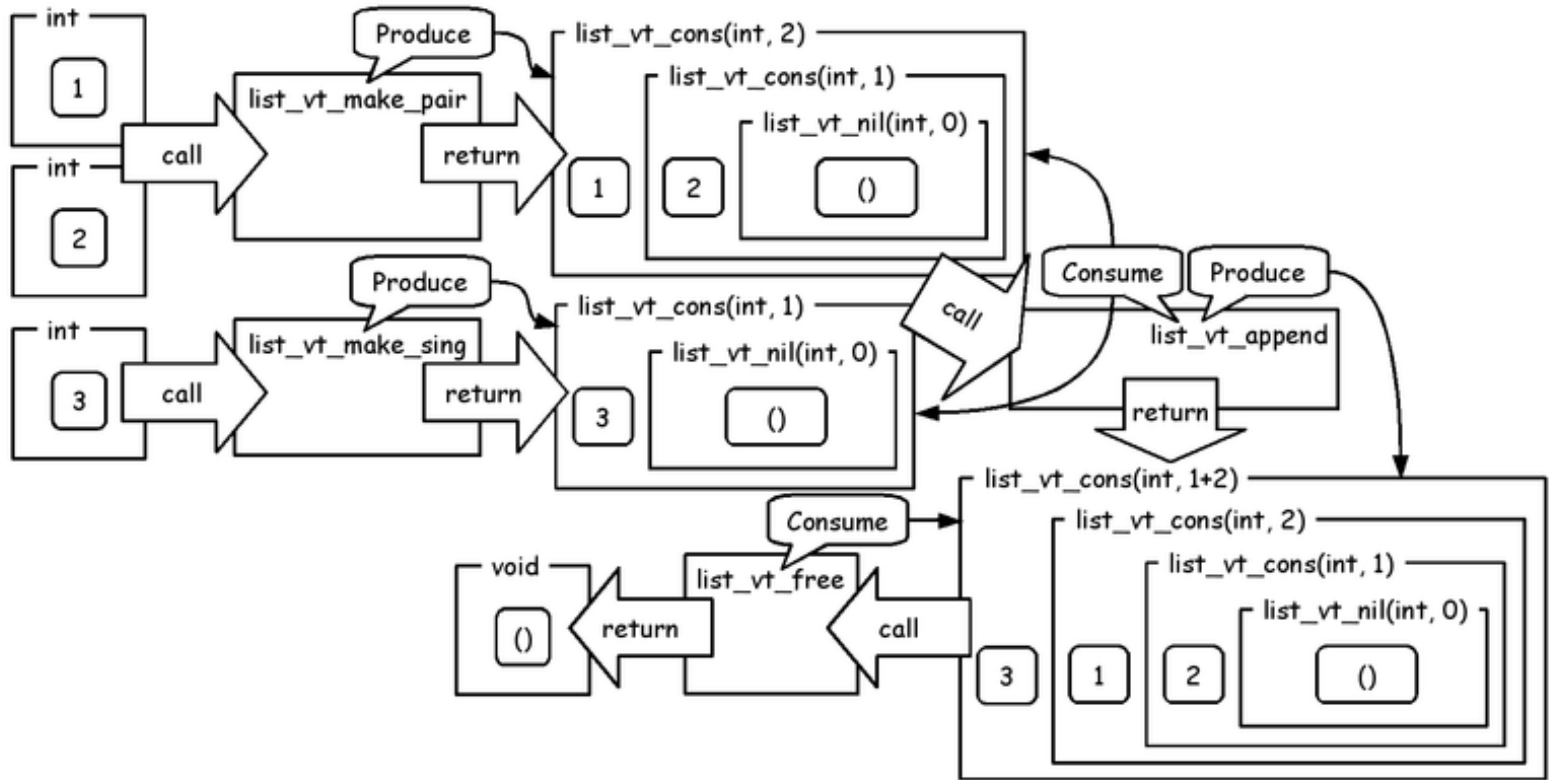
primplment gds = Gd_Sun ()



[4] View (linear type) on ATS

- ☆ View is linear type in ATS
- ☆ View manages producing and consuming resource
- ☆ Example of resource: memory chunk, array, list, queue, lock/unlock, session, ...
- ☆ At-view is a ticket to permit dereferencing pointer

At-view: Figure of append list



At-view: Code of append list

```
#include "share/atspre_staload.hats"
```

```
implement main0 () = {  
  val l1 = list_vt_make_pair (1, 2)  
  val l2 = list_vt_make_sing 3  
  val l3 = list_vt_append (l2, l1)  
  val () = print_list_vt<int> l3 (* => 3, 1, 2 *)  
  (*val () = list_vt_free<int> l3 *) (* Compile-time error! *)  
}
```

[5] Compare ATS with Event-B

- ☆ Event-B generates code from verification.
- ☆ ATS write verification in code.
- ☆ Event-B focuses on interaction between contexts.
- ☆ ATS focuses on specification in single context.
- ☆ ATS can manage multi-context with some code.

Compare ATS with VeriFast

VeriFast is very similar to ATS, but...

- ☆ VeriFast is closed source.
- ☆ ATS is open source.
- ☆ VeriFast is for C/Java.
- ☆ ATS is for C/Erlang/JavaScript/PHP/Perl/Python.
- ☆ VeriFast's host language is C or Java.



Compare ATS with XMOS

☆ XMOS strongly focuses on channel implemented on hardware.

☆ ATS also uses strongly typed channel what needs some support by OS.

<http://ats-lang.sourceforge.net/EXAMPLE/EFFECTIVATS/ssntyped-channels-1/main.html>

☆ ATS's channel may be easily used without OS-support, with hardware-support like XMOS.

[6] Conclusion

- ☆ ATS can use strong type without any OS.
- ☆ ATS can prove code using dependent types.
- ☆ ATS can safely use pointer using linear types.
- ☆ ATS can manage resource using linear types.
- ☆ Join "ATS logo Japan ATS User Group"!
- ☆ <http://jats-ug.metasepi.org/>